## Problem sheet 1

## 1. Lagrange interpolation formula

Let $x_{0}, x_{1}, \ldots, x_{n}$ be $(n+1)$ pairwise distinct points and let there be given $(n+1)$ arbitrary numbers $y_{0}, y_{1}, \ldots, y_{n}$. Further define the fundamental Lagrange polynomials by

$$
l_{i}^{n}(x)=\frac{\left(x-x_{0}\right) \ldots\left(x-x_{i-1}\right)\left(x-x_{i+1}\right) \ldots\left(x-x_{n}\right)}{\left(x_{i}-x_{0}\right) \ldots\left(x_{i}-x_{i-1}\right)\left(x_{i}-x_{i+1}\right) \ldots\left(x_{i}-x_{n}\right)} .
$$

(a) Show that

$$
p_{n}(x)=\sum_{i=0}^{n} y_{i} l_{i}^{n}(x)
$$

is the unique polynomial $p_{n}(x)$ of degree $n$ such that

$$
p_{n}\left(x_{i}\right)=y_{i}, \quad \text { for } i=0,1, \ldots, n \text {. }
$$

Hint: You may first show that $l_{i}^{n}\left(x_{j}\right)$ equals 1 if $i=j$ and zero if $i \neq j$.
(b) Let $\Pi(x)=\left(x-x_{0}\right)\left(x-x_{1}\right) \ldots\left(x-x_{n}\right)$ and show that

$$
p_{(2 n+1)}(x)=\sum_{i=0}^{n}\left[f\left(x_{i}\right)\left(1-\frac{\Pi^{\prime \prime}\left(x_{i}\right)}{\Pi^{\prime}\left(x_{i}\right)}\left(x-x_{i}\right)\right)+f^{\prime}\left(x_{i}\right)\left(x-x_{i}\right)\right]\left(l_{i}^{n}(x)\right)^{2}
$$

is the unique polynomial of degree $(2 n+1)$ for which

$$
p_{(2 n+1)}\left(x_{i}\right)=f\left(x_{i}\right), \quad p_{(2 n+1)}^{\prime}\left(x_{i}\right)=f^{\prime}\left(x_{i}\right), \quad \text { for } i=0,1, \ldots, n
$$

## 2. Orthogonal polynomials

Let $P_{0}(x), P_{1}(x), \ldots$ be a set of orthonormal polynomials, e.g.

$$
\int_{a}^{b} d x w(x) P_{i}(x) P_{j}(x)=\delta_{i j}
$$

let $x_{1}, x_{2}, \ldots, x_{n+1}$ be the zeros of $P_{n+1}(x)$ and $w_{1}, w_{2}, \ldots, w_{n+1}$ the corresponding Gaussian weights given by

$$
w_{j}=\int_{a}^{b} d x w(x) l_{j}^{n+1}(x)
$$

Show that for $i, j<n+1$

$$
\sum_{k=1}^{n+1} w_{k} P_{i}\left(x_{k}\right) P_{j}\left(x_{k}\right)=\delta_{i j}
$$

e.g. the $P_{0}, P_{1}, \ldots, P_{n}$ are orthonormal on the zeros of $P_{n+1}$. This equation can be useful to check the accuracy with which the zeros and weights of $P_{n+1}$ have been determined numerically.

## 3. Monte Carlo integration

The Monte Carlo estimate for the $d$-dimensional integral

$$
I=\int d x f(x)=\int d^{d} u f\left(u_{1}, \ldots, u_{d}\right)
$$

is given by

$$
E=\frac{1}{N} \sum_{n=1}^{N} f\left(x_{n}\right) .
$$

In order to discuss the error estimate for finite $N$, we first introduce the variance $\sigma^{2}(f)$ of the function $f(x)$ :

$$
\sigma^{2}(f)=\int d x(f(x)-I)^{2}
$$

Show that

$$
\int d x_{1} \ldots \int d x_{N}\left(\frac{1}{N} \sum_{n=1}^{N} f\left(x_{n}\right)-I\right)^{2}=\frac{\sigma^{2}(f)}{N},
$$

e.g. the variance with which the Monte Carlo estimate differs from the true value is given by $\sigma^{2}(f) / N$.
Hint: You may want to introduce an auxiliary function $g(x)=f(x)-I$ and show $\int d x g(x)=0$ first.

## 4. Pseudo-random numbers

Pseudo-random numbers were in the early days often generated by a multiplicative linear congruential generator according to the prescription

$$
s_{i}=\left(a s_{i-1}+c\right) \bmod m
$$

Consider the simple multiplicative linear congruential generator with $a=5, c=1, m=16$ and $s_{0}=1$. Write down the first twenty numbers generated with this method. How long is the period? Write down the sequence also in the binary representation and look at the lowest order bits.

## 5. Gray code

The Gray code representation for a number $n$ is a binary representation such that the representations for $n$ and $n+1$ differ in only one bit. The Gray code representation can be obtained from the binary representation according to

$$
\ldots g_{3} g_{2} g_{1}=\ldots b_{3} b_{2} b_{1} \oplus \ldots b_{4} b_{3} b_{2}
$$

where $\oplus$ deontes the bitwise exclusive-or operation ( $0 \oplus 0=1 \oplus 1=0,0 \oplus 1=1 \oplus 0=1$ ). Find the Gray code representation for $n=0,1, \ldots, 7$.

